



Loading and Saving

- ◆ Interpreter
- ◆ First prompt (-) and secondary prompt (=)
- ◆ Need ; after each expression/definition
- ◆ Loading ML source text from a file
 - Create a file named "myfile.sml"
 - Either start ML
 - and use the function `use: string -> unit`
 - `use "c:\\myfile.sml";`
 - Don't forget the double '\\' in the path!
 - Or redirect the input and output
 - `C:\ sml < myfile.sml > output`

ML Introduction.3

A simple tutorial

- ◆ ML is usually used as interpreter (a compiler is also available)
- ◆ Expressions followed by a semicolon yield a response
 - `2+2;`
 - `val it = 4 : int`
- ◆ Doing simple arithmetic
 - `3.2 - 2.3;`
 - `val it = 0.9 : real`
 - `Math.sqrt(2.0);`
 - `val it = 1.414213562 : real`

ML Introduction.4

Declaring Constants

- ◆ Naming constants
 - `val seconds = 60;`
`val seconds = 60 : int`
- ◆ minutes per hour
 - `val minutes = 60;`
`val minutes = 60 : int`
- ◆ hours per day
 - `val hours = 24;`
`val hours = 24 : int`
- ◆ Using names in expressions
 - `seconds * minutes * hours;`
`val it = 86400 : int`

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The identifier '*it*'

- ◆ By referring to *it*, one can use the last value
 - `it div 24;`
`val it = 3600 : int`
- ◆ Any previous value of *it* is lost unless saved
 - `val secsinhour = it;`
`val secsinhour = 3600 : int;`
- ◆ Underscores can be used in names for readability
 - `val secs_in_hour = seconds * minutes;`
`val secs_in_hour = 3600 : int`

ML Introduction.6

Legal Names - Alphabetic Names

- ◆ Alphabetic name
 - Begins with a letter
 - Then followed by letters, digits, underscore, or single quotes
 - Examples:
 - `x`
 - `UB40`
 - `Hamlet_Prince_of_Denmark`
 - `h''3_H`
 - `or_any_other_name_that_is_as_long_as_you_like`
 - Case of letters matters
 - ML was designed by Mathematicians who like primes
 - `x`, `x'`, `x''`, `x'''`

ML Introduction.7

Legal Names - Symbolic Names

- ◆ Permitted over the characters:
`! % & $ # + - * / : < = > ? @ \ ~ ` ^ |`
- ◆ May be as long as you like:
`---->`
`$^$^$^$`
`!!?@**??!!`
`:_||==>->#`
- ◆ Should not be one of the ML reserved special syntax:
`: _ | = => -> #`
- ◆ Allowed whenever an alphabetic name is:
`- val +--+ = 1415;`
`val +--+ = 1415 : int`

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ML's keywords

`abstype` `and` `andalso` `as` `case` `datatype` `do` `else`
`end` `eqtype` `exception` `fn` `fun` `functor` `handle` `if`
`in` `include` `infix` `infixr` `let` `local` `nonfix` `of` `op`
`open` `orelse` `raise` `rec` `sharing` `sig` `signature`
`struct` `structure` `then` `type` `val` `while` `with`
`withtype`

- Avoid ML's keywords when choosing names
- Especially watch out from the short ones:
`as` `fn` `if` `in` `of` `op`

ML Introduction.9

ML Primitive Types

`int`, `real`, `string`, `char`, `bool`, `unit`

ML Introduction.10

Integer Types

- ◆ Constants
 - sequence of digits
 - 0
 - 01234
 - ~ for a unary minus sign
 - ~23
 - ~85601435654638
- ◆ Infix operations:
 - + - * div mod
- ◆ Conventional precedence
 - $((m * n) * 1) - (m \text{ div } j) + j$
 - parenthesis can be dropped without change of meaning.

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Real types

- ◆ Constants
 - decimal point
 - 0.01
 - 2.718281828
 - E notation
 - 7E~5
 - ~1.2E12
 - ~123.4E~2 is the same as ~1.234
 - ~ for unary minus
- ◆ Infix operators
 - + - * /
- ◆ Functions
 - `floor(r)` converts real to int, `real(i)` converts int to real
 - `sqrt`, `sin`, `cos`, `tan`, `exp`, `ln` all of type *real* -> *real*
 - All need the **Math** prefix: `Math.sqrt`, `Math.sin`
 - Infix operators have lower precedence.

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Strings

- ◆ Constants are written in double quotes
 - `"ML is the best";`
 - `val it = "ML is the best" : string`
- ◆ Special characters `\n` `\t` `\"` `\\`
- ◆ Concatenation
 - `"Standard" ^ " ML";`
 - `val it = "Standard ML" : string`
- ◆ `size` returns the number of characters
 - `size (it);`
 - `val it = 11 : int`
- ◆ Infix operators
`<`, `^`

`size("") is 0`

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Characters

- ◆ Chars are distinguished from strings of length 1 by the # sign
 - `"0";`
 - `val it = "0" : string`
 - `#"0";`
 - `val it = #"0" : char`
- ◆ Converting between strings and characters using `str` and `sub`
 - `str("0");`
 - `val it = "0" : string`
 - `String.sub("hello",0);`
 - `val it = #"h" : char`
- ◆ Converting chars to and from ASCII using `ord` and `chr`
 - `ord #"0";`
 - `val it = 48 : int`
 - `chr it;`
 - `val it = #"0" : char`

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Boolean

- ◆ The two values are
 - `true`;
 - `val it = true : bool`
 - `false`;
 - `val it = false : bool`

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Tuples Cartesian Product Type

- ◆ `(x1, x2, ..., xn)`
 - The n-tuple whose components are `x1, x2, ..., xn`.
 - The components can be of any type, including tuples.
- ◆ Examples
 - `val a = (1.5, 6.8);`
 - `val a = (1.5, 6.8) : real * real`
 - `(1, 1.5);`
 - `val it = (1, 1.5) : int * real`
 - `("str", 1, true, (#"0", 0.1));`
 - `val it = ("str", 1, true, (#"0", 0.1)) : string * int * bool * (char * real)`

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Records

- ◆ Records have components (fields) identified by name
 - `val me = { name="Ofir", age=30 };`
 - `val me = {age=30,name="Ofir"} :`
`{age:int, name:string}`
- ◆ Type lists each field as `label : type`
- ◆ Enclosed in braces `{ ... }`
- ◆ Selecting a field
 - `#name(me);`
 - `val it = "Ofir" : string`
- ◆ Tuples can be seen as records with numbers as implicit field labels
 - (x_1, x_2, \dots, x_n) is $\{1=x_1, 2=x_2, \dots, n=x_n\}$
 - `#2 ("one", "two", "three");`
 - `val it = "two" : string`

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Lists

- ◆ A list is a finite sequence of elements.
 - `[3,5,9]`
 - `["a", "list"]`
 - `[]`
- ◆ Elements may appear more than once
 - `[3,4]`
 - `[4,3]`
 - `[3,4,3]`
 - `[3,3,4]`
- ◆ Elements may have any type. But all elements of a list must have the same type.
 - `[(1,"One"),(2,"Two")] : (int*string) list`
 - `[[3.1],[],[5.7, -0.6]]: (real list) list`

ML Introduction.18

Mapping - Functions

- ◆ - `fun sq(x:int) = x*x;`
`val sq = fn : int -> int`
 - keyword `fun` starts the function declaration
 - `sq` is the function name
 - `x:int` is the formal parameter with type constraint
 - `x*x` is the body and it is an **expression**
 - the type of a function is printed as `fn`
 - The result of the function is the result of evaluating the **expression** of the function body with the actual parameter
 - `int -> int` is the standard mathematical notation for a function type that takes a real number and returns a real number

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Applying a Function

- ◆ Simple function call
 - `sq (3);`
`val it = 9 : int`
- ◆ When a function is called the parameter is evaluated and then passed to the function (seems obvious but it is not always the case in functional languages...)
 - `sq (sq(3));`
`val it = 81 : int`
- ◆ The parentheses around the argument are optional
 - `sq 3;`
`val it = 9 : int`
- ◆ Parentheses are also optional in function definitions
 - `fun sq x:int = x*x;`
`val sq = fn: int -> int`

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Arguments and Results

- ◆ Every function has one argument and one result.
- ◆ Any type can be passed/returned !!!
- ◆ Tuples are used to pass/return several arguments
 - `val a = (1.5, 6.8);`
 - `val a = (1.5, 6.8) : real * real`
 - `fun lengthvec (x:real,y:real) = sqrt(x*x + y*y);`
 - `val lengthvec = fn : real * real -> real`
 - `lengthvec a;`
 - `val it = 6.963476143 : real`
 - `fun negvec (x:real,y:real) = (~x, ~y);`
 - `val negvec = fn : real * real -> real * real`
 - `negvec (1.0, 1.0);`
 - `val it = (~1.0, ~1.0) : real * real`

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Functions as Values

- ◆ Anonymous functions with `fn` notation
 - `fn x => x*x;`
 - `val it = fn : int -> int`
 - `it(3);`
 - `val it = 9 : int`
- ◆ The following declarations are identical
 - `fun sq x = x*x;`
 - `val sq = fn x => x*x`

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Functions as Parameters

- ◆ The definition of sigma:

$$\sum_{i=x}^y f(i) = \begin{cases} f(x) + \sum_{i=x+1}^y f(i) & \text{if } x \leq y \\ 0 & \text{otherwise} \end{cases}$$

- ◆ Functions can be given as parameters to other functions

```
- fun Sigma(f,x,y) =  
=   if x<=y then f(x) + Sigma(f,x+1,y)  
=   else 0;  
val Sigma =  
    fn : (int -> int) * int * int -> int  
- Sigma(sq,1,3);  
val it = 14 : int  
- Sigma(fn x => x*x,1,3);  
val it = 14 : int
```

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Function as Return Value

- ◆ Functions can also be **returned** from other functions

```
- fun inttwice(f:(int->int)) =  
    fn x => f(f(x));  
val inttwice = fn : (int -> int) -> int -> int
```

- ◆ The arrow associates to the right so the last line is equivalent to

```
val inttwice = fn : (int -> int) -> (int -> int)
```

- ◆ Example

```
- inttwice(fn x => x*x);  
val it = fn : int -> int  
- it(3);  
val it = 81 : int
```

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Type Inference

- ◆ ML deduces the types in expressions
- ◆ Type checking the function:

```
fun facti (n,p) =  
  if n=0 then p else facti(n-1,n*p);
```

- constants `0` and `1` have type `int`
- therefore `n=0` and `n-1` involve integers
- so `n` has type `int`
- `n*p` must be integer multiplication, so `p` has type `int`
- `facti` returns type `int`, and its argument type is `int*int`

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Type Constraints

- ◆ Certain functions are overloaded, e.g., `abs`, `+`, `-`, `~`, `*`, `<`.
- ◆ Type of an overloaded function is determined from context, or is set to `int` by default.
- ◆ Types can be stated explicitly.
- ◆ Examples:

```
- fun min(x,y) = if x < y then x else y;  
val min = fn : int * int -> int  
- fun min(x:real,y) = if x < y then x else y;  
val min = fn : real * real -> real  
- fun min(x:string,y) = if x < y then x else y;  
val min = fn : string * string -> string  
- fun min(x,y):real = if x < y then x else y;  
val min = fn : real * real -> real  
- fun min(x,y) = if x < y then x:real else y;  
val min = fn : real * real -> real
```

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Polymorphic type checking

- ◆ Weakly typed languages (e.g.. Lisp)
 - give freedom
- ◆ Strongly typed languages (e.g. Pascal)
 - give security by restricting the freedom to make mistakes
- ◆ Polymorphic type checking in ML
 - security of strong type checking
 - great flexibility (like weak type checking)
 - most type information is deduced automatically
 - an object is polymorphic if it can be regarded as having any kind of type

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Polymorphic function definitions

- ◆ If type inference leaves some types completely unconstrained then the definition is polymorphic
- ◆ A polymorphic type contains a *type variable*, e.g. 'a
- ◆ Example:

```
- fun pairself x = (x,x);  
val pairself = fn : 'a -> 'a * 'a  
  
- pairself 4.0;  
val it = (4.0,4.0) : real * real  
  
- pairself "NN";  
val it = ("NN","NN") : string * string  
  
- pairself (1.0,3);  
val it = ((1.0,3),(1.0,3)):(real*int)*(real*int)  
  
- fun pair (x,y) = (y,x);  
val pair = fn: ('a * 'b) -> ('b * 'a)
```

ML Introduction.28

Functions as Values The Polymorphic Case

```
- fun twice f = fn x => f(f(x));  
val twice = fn : ('a -> 'a) -> 'a -> 'a  
- fun ident x = x;  
val ident = fn : 'a -> 'a  
- twice (fn x => x*x);  
val it = fn : int -> int  
- it(2);  
val it = 16 : int
```

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Functions as Values The Polymorphic Case (cont.)

- ◆ Sometimes ML gives us hard time when we give polymorphic value to polymorphic function. For example:

```
- twice ident;  
stdIn:... Warning: type vars not generalized  
because of value restriction are  
instantiated to dummy types (X1,X2,...)  
val it = fn : ?.X1 -> ?.X1
```
- ◆ The reason for this is outside the scope of this course. You usually may ignore it. Or, if needed, workaround the problem:

```
- fn x => (twice ident)(x);  
val it = fn : 'a -> 'a
```

ML Introduction.30

Functional vs. Imperative

- ◆ Imperative - using commands to change the state.
- ◆ Functional - stateless. Using expressions recursively to calculate the result.
- ◆ Example: Euclid's Algorithm for the Greatest Common Divisor (GCD) of two natural numbers:

$$\text{gcd}(m, n) = \begin{cases} n & m = 0 \\ \text{gcd}(n \bmod m, m) & m > 0 \end{cases}$$

How would a GCD program look like in functional vs. imperative language?

ML Introduction.31

GCD - Pascal vs. ML

- ◆ An imperative Pascal Program:

```
function gcd(m,n: integer): integer;  
var prevm: integer;  
begin  
  while m<>0 do begin  
    prevm := m; m := n mod m; n := prevm  
  end;  
  gcd := n  
end;
```

- ◆ A functional program in Standard ML:

```
fun gcd(m,n) =  
  if m=0 then n else gcd(n mod m, m);
```

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